

A Consolidated Authentication Model in Cloud Computing Environments

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Abstract

Due to increasing needs of Internet access through smart phones and smart pads, it is essential to have service provider systems, which allows to access services through a variety of devices. In particular, this system is required to protect credential and personal information saved in each device, is need a more efficient and secure consolidated authentication model (CAM) in order to authenticate a user and devices.

This paper analyzes the current user authentication model for both user and device authentication and securely available credential (SACRED) standards. Furthermore, it is also our intention to design N-screen based consolidated user authentication model that meets framework and protocol requirement of credentials and privacy protection requirements in a cloud computing environments.

Keywords: *User Authentication, Privacy, PKI, Credentials, N-screen*

1. Introduction

At each site, to deracinate a disclosure, abuse, and misuse of personal information, Privacy Act was enacted and enforced. [1] In addition, although accredited certificate is the most commonly used method for user authentication [2] in various e-governments and online based financial sector, the risk of using it has been coming on the rise due to spreading of malicious program such as virus, spyware or malware. In order to solve this problem, there has been active discussion over ways to come up with enhanced safety measures. In this regard, it is necessary to develop new technologies which enable the accredited certificate to utilize in various smart devices and cloud computing environment. [3]

In this paper, it is our purpose to suggest a safe and convenient user authentication model that mobile device users can effortlessly use credentials in cloud computing environments. After this brief introduction, the reminder of this paper is organized as follows: section 2 discusses the security issues in clouds and M2M (Machine to Machine) environments, the analysis of current user authentication model, and securely available credentials. Then we propose consolidated user authentication with framework architecture and protocol framework is shown in section 3. Section 4 describes overall prototyping of proposed model. Section 5 presents the comparison and verification of our architecture. Finally we conclude the paper and suggest the future research work in section 6.

In this user authentication model, credentials are stored in a specific smart device to generate a digital signature. However, since cloud computing environments should support a variety of devices, a more efficient and secure user authentication model is required. In addition to that, tight security system needs to be equipped in dealing with the life cycle of personal information which ranges from collection, storage, use, transfer to disposal.

2.3. Securely available credentials

The SACRED Working Group is working on the standardization of a set of protocol for securely transferring credentials among devices. The international standards of SACRED Working group consist of RFC 3157(Requirements) [9], RFC 3760 (Credential Server Framework) [10], and RFC 3767(Protocol) [11].

Problems and limitations of existing SACRED standards are as below.

First, although it defines the framework and protocol requirements with respect to securely available credential, the existing SACRED standards does not provide a detail implementation guideline. Second, SACARE defines the upload and download protocol for credentials, but does not define a protocol to create a proxy signature from Signing Server using uploaded credential by Client.

To solve these problems, let us redefine a credential framework that meets the framework and protocol requirements of SACRED and design a credential protocol based on ASN.1. In addition, we would like to define a protocol of credential roaming and proxy signature which fits in cloud computing environments.

3. Consolidated Authentication Model (CAM)

3.1. Overview

The CAM consists of consolidated authentication mechanism and policy compliance mechanism. The consolidated authentication mechanism is guaranteed to the use's consolidated authentication using security technology. The policy compliance mechanism is supported the systemically policy not only managing but also controlling the system during interoperation process.

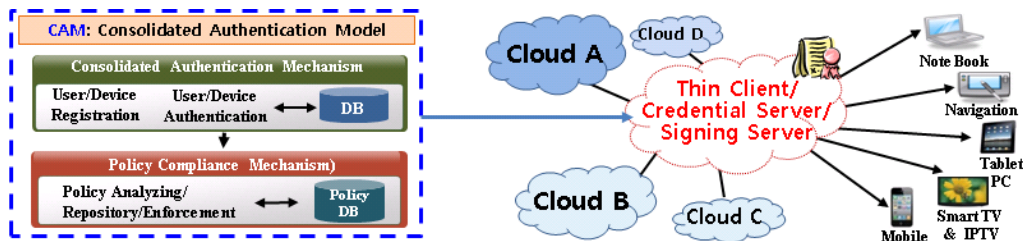


Figure 2. CAM Architecture

3.2. Framework architecture

Client, Credential Server, Signing Server in framework perform the following roles. A Client uploads or downloads credentials from Credential Server through a variety of devices such as PC, smart pad, and smart phone and generates the digital signature from Signing Server. The Credential Server (CS) downloads secure credentials and uploads

them from the client. The Credential Store is the repository for secured credentials. The Signing Server (SS) creates a digital signature by the Client's request.

3.3. Secure credentials design

Credentials are information that can be used to establish the identity of an entity, or help that entity communicate securely. Credentials include such things as private keys, trust roots, tickets, or the private part of a Personal Security Environment (PSE) [RFC2510]. [9] Several standardized formats for the representation of credentials exist e.g., PKCS#12[12], PKCS#15[13]. Secure Credentials is a set of one or more credentials that have been cryptographically secured, e.g., encrypted/MACed with a passkey [11].

3.4. Protocol framework

Consolidated authentication mechanism consists of account management module (AMM), credential roaming module (CRM), and proxy signature module (PSM).

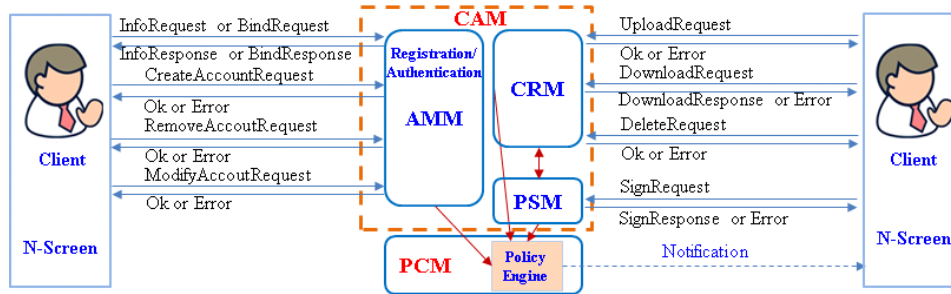


Figure 3. Protocol Framework

Notations and abbreviation for protocol design is as follows:

Table 1. Notations and Abbreviation

Symbol: Description	Symbol: Description	Symbol: Description
SN: Serial Number	TIME: Signed Time	KDF2: Key Derivation Function 2[14]
SD: Signed Data	PV: Password Verifier	RC, RS: Random Number
ID: Identification Number	C: Credential	SC: Secure Credential, K: Key for SC
H(): Hash Function	E(): Encryption	S(): Generate Signature
=?: Compare with	D(): Decryption	V(): Verify Signature

3.4.1. Initialization and key sharing operations

1) BindRequest/BindResponse Protocol

If Client has his/her own digital certificate, he/she creates a digital signature using their private key and passes it to credential server.

$$① \text{ Client: } I=H(ID), N=H(SN), SD=S_{client_key}(I|N|TIME)$$

$$② \text{ Client } \rightarrow \text{ CS: BindRequest}(SD, Client_Cert)$$

$$③ \text{ CS: } I|N|TIME=V_{client_cert}(SD), ERS=E_{client_cert}(RS)$$

$$④ \text{ CS } \rightarrow \text{ Client: BindResponse}(ERS, CS_cert) \quad ⑤ \text{ Client: } RS=D_{client_key}(ERS)$$

2) InfoRequest/InfoResponse Protocol

If Client doesn't have the certificate, Credential Server sends a session key (RS) after the server creates secure channel through SSL/TLS [15] or DH key exchange [16].

- ① Client → CS: InfoRequest ② CS → Client: InfoResponse(RS, CS_cert)

3.4.2. Account management operations

1) Create Account Protocol

When Client creates user account, Credential Server registers hash value which contains both client's unique ID such as Resident Registration number and unique device information such as serial number or MAC address that is used by the Client PC, smart pad, smart phone, etc.

- ① Client → CS: BindRequest(SD, Client_Cert) ② CS → Client: BindResponse(ERS, CS_cert)
 ③ Client: $\underline{ERC} = E_{RS}(RC)$, $I = H(ID)$, $N = H(SN)$, $K = KDF2(I, PW)$, $PV = H(H(ID, K))$,
 $HI = H(I|PV|RS|RC|N)$, $\underline{EI} = E_{RC}(I|PV|HI|N)$
 ④ Client → CS: CreateAccountRequest(ERC, EI)
 ⑤ CS: $RC = D_{RS}(ERC)$, $I|PV|HI|SN = D_{RC}(EI)$, $HI' = H(I|PV|RS|RC|SN)$, $I|PV|HI|SN = ?I|PV|HI|SN$
 ⑥ CS → Client: CreateAccountResponse(Ok or Error)

2) Modify Account Protocol and Remove Account Protocol

Client can register new device information or modify the registered device information. If the Client no longer uses the account, a registered account can be removed.

- ① Client → CS: BindRequest(SD, Client_Cert) ② CS → Client: BindResponse(ERS, CS_cert)
 ③ Client: $\underline{ERC} = E_{RS}(RC)$, $I = H(ID)$, $N = H(SN)$, $K = KDF2(I, PW)$, $PV' = H(ID, K)$,
 $HI = H(I|PV'|RS|RC|N)$, $\underline{EI} = E_{RC}(I|PV'|HI|N)$
 ④ Client → CS: {Modify, Remove}AccountRequest(ERC, EI)
 ⑤ CS: $RC = D_{RS}(ERC)$, $I|PV'|HI|N = D_{RC}(EI)$, $HI' = H(I|PV'|RS|RC|N)$,
 $I|H(PV')|HI|N = ?I|PV|HI|N$
 ⑥ CS → Client: {Modify, Remove}AccountResponse (Ok or Error)

3.4.3. Credential roaming operations

1) Credential Upload Protocol

The registration process that Client uploads credentials to Credential Server is as below.

- ① Client → CS: BindRequest(SD, Client_Cert) ② CS → Client: BindResponse(ERS, CS_cert)
 ③ Client: $\underline{ERC} = E_{RS}(RC)$, $I = H(ID)$, $N = H(SN)$, $K = KDF2(I, PW)$, $PV = H(H(ID, K))$,
 $SC = E_K(C)$, $HI = H(I|PV|SC|RS|RC|N)$, $\underline{EI} = E_{RC}(I|PV|SC|HI|N)$
 ④ Client → CS: UploadRequest(ERC, EI)
 ⑤ CS: $RC = D_{RS}(ERC)$, $I|PV|SC|HI|N = D_{RC}(EI)$, $HI' = H(I|PV|SC|RS|RC|N)$,
 $I|PV|SC|HI|N = ?I|PV|SC|HI|N$ ⑥ CS → Client: UploadResponse(Ok or Error)

2) Credential Download Protocol from Credential Server

In order to use the credential in a variety of environments, the download procedure of credentials from the credential server is as follows.

- ① Client → CS: InfoRequest ② CS → Client: InfoResponse(RS, CS_cert)
 ③ Client: $\underline{ERC} = E_{RS}(RC)$, $I = H(ID)$, $N = H(SN)$, $K = KDF2(I, PW)$, $PV' = H(ID, K)$,
 $HI = H(I|PV'|RS|RC|N)$, $\underline{EI} = E_{RC}(I|PV'|HI|N)$
 ④ Client → CS: DownloadRequest(ERC, EI)
 ⑤ CS: $RC = D_{RS}(ERC)$, $I|PV'|HI|N = D_{RC}(EI)$, $HI' = H(I|PV|RS|RC|N)$,
 $I|H(PV')|HI|N = ?I|PV|HI|N$, $\underline{ESC} = E_{RC}(SC)$
 ⑥ CS → Client: DownloadResponse(ESC) ⑦ Client: $SC = D_{RC}(ESC)$, $C = D_K(SC)$

3) Credential Download Protocol from direct solutions

The way to deliver credential among different devices is through PKCS#12, which is currently supported by most browsers. Credential is double-protected by the password of private key and that of PKCS#12.

- ① Device 1: $PKCS\#12Export(data\ or\ file(*.pfx\ or\ *.p12))$
- ② Device 1 \rightarrow Device 2: Transfer PKCS#12 data
- ③ Device 2: $PKCS\#12Import(data\ or\ file(*.pfx\ or\ *.p12))$

3.4.4. Proxy signature operations

1) Proxy Signature [17] Protocol

The signing process of Client using Signing Server is as follows:

- ① $SS \rightarrow CS: BindRequest(SD, SS_cert)$ ② $CS \rightarrow SS: BindResponse(ERS, CS_cert)$
- ③ $Client \rightarrow SS: InfoRequest$ ④ $SS \rightarrow Client: InfoResponse(RS, SS_cert)$
- ⑤ $Client: \underline{ERC} = E_{RS}(RC), I = H(ID), N = H(SN), K = KDF2(I, PW), PV' = H(ID, K), \underline{D} = H(M),$
 $HI = H(I|PV'|D|RS|RC|N), EI = E_{RC}(I|PV'|HI|N), \underline{ED} = E_{RS}(ERC/EI), \underline{K} = E_{RS}(K)$
- ⑥ $Client \rightarrow SS: SignRequest(ED, D, EK, ERC)$ ⑦ $SS: ERC|EI = D_{RS}(ED)$
- ⑧ $SS \rightarrow CS: DownloadRequest(ERC, EI)$
- ⑨ $CS: RC = D_{RS}(ERC), I|PV'|HI|N = D_{RC}(EI), HI' = H(I|PV'|RS|RC|N),$
 $I|H(PV')|HI|N = ? I|PV'|HI|N, \underline{ESC} = E_{RC}(SC)$ ⑩ $CS \rightarrow SS: DownloadResponse(ESC)$
- ⑪ $SS: RC = D_{RS}(ERC), SC = D_{RC}(ESC), K = D_{RS}(EK), C = D_K(SC), \underline{SD} = S_{client_key}(D)$
- ⑫ $SS \rightarrow Client: SignResponse(SD)$ ⑬ $Client: D = H(M), D' = V_{client_cert}(SD), D = ? D'$

4. Prototyping

4.1. Implementation

To demonstrate the feasibility of our architecture, we implemented a prototype system which provides consolidate user authentication for secure system. This system is developed using JSP, JAVA, iPhone development toolkit technologies. [18] This table below shows CAM's add account, signing procedure using iPhone's Application and user certificate saved in server.



Figure 4. iOS User Interface of CAM

4.2. Simulation

The simulation platform and certificate to show performance difference between CAM and conventional PKI is as follows;

Table 2. Simulation Environments

Category	Description
Platform	PC (Intel dual core, 3.2GHz), Smart Phone (iPhone 4), UNIX (SUN Fire V240 1.5GHz*2ea)
certificate	User certificate, Device certificate: RSA2048bit/SHA256

A general scenario of CAM and PKI is as follows:

1) CAM: User's device requests proxy signature to Signing server and service provider verify that result. 2) PKI: User's device generates digital signature and service provider verify it. We simulated RSA signing and verification for each platform. The CAM provide more a time saving and enhanced security than traditional authentication using PKI by providing consolidated authentication for user, device and contents.

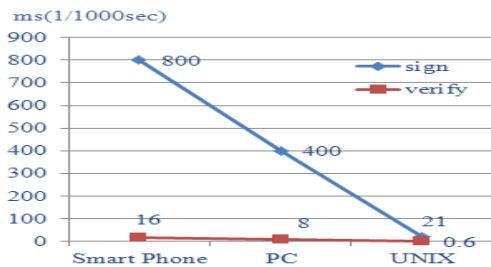


Figure 5. Result of RSA signing and verify

Table 3. Simulation Result

Category	PKI	CAM
User authentication	Slow (PC, Smart Phone)	Fast (UNIX)
Authentication method	Provided by each service	Centralized management

5. Comparison and verification

We demonstrate that CAM architecture can solves the existing problems by satisfying framework requirements, protocol requirements, privacy protection requirements, and by comparing with the current user authentication model.

The CAM satisfies framework requirements as follow:

Table 4. Framework Requirements

No	Description	CAM
F1	The framework must support both "credential server" and "direct" solutions.	O
F2	The "credential server" and "direct" solutions should use the same technology.	O
F3	The framework must allow for protocols which support different user authentication schemes.	O
F4	The details of the actual credential type or format must be opaque to the protocol.	O
F5	The framework must allow use of different transports.	O

The CAM satisfies protocol requirements as follow:

Table 5. Protocol Requirements

No	Description	CAM
G1	Credential transfer both to and from a device must be supported.	O
G2	Credentials must not be forced by the protocol to be present in cleartext at any device other than the end user's.	O
G3	The protocol should ensure that all transferred credentials be authenticated in some way.	O
G4	The protocol must support a range of cryptographic algorithms.	O
G5	The protocol must allow the use of various credential types and formats.	O
G6	One mandatory to support credential format must be defined.	O
G7	One mandatory to support user authentication scheme must be defined.	O
G8	The protocol may allow credentials to be labeled with a text handle.	O
G9	Full I18N support is required (via UTF8 support)	O
G10	The protocol is able to support privacy, that is, anonymity for the client.	O
G11	Transferred credentials may incorporate timing information.	O

The CAM compare with the old user authentication model as follow:

Table 6. Comparison between Old model and CAM

Category	Framework requirements	Protocol requirements	Privacy protection requirements	Proxy signature
Old model	Δ	Δ	Δ	X
CAM	O	O	O	O

O: provided Δ: partially provided X: not provided

6. Conclusion and Future Work

In this paper, we discuss the security and privacy issues of the current user authentication model that are not able to provide credential roaming in cloud computing environments due to the absence of securely available credential protocol in consolidated user authentication method. In order to solve this problem, we proposed the secure CAM architecture so that one credential is applicable to various mobile devices in cloud computing environments.

Following is contributions of N-screen based consolidated user authentication model for internet services that meets framework and protocol requirement of credentials and privacy protection requirements in a cloud computing environments. We designed the secure CAM architecture in cloud computing environments, which not only provides more flexible authentication framework but also leads to safer credential management in operating various mobile devices such as smart phone, smart pad, etc. We define framework architecture, credential profile, protocol framework for consolidated authentication mechanism in order to provide an appropriate user authentication model for a cloud computing environments.

The future study will continue to focus the design and implement of our suggested model, and we will expand to new devices and environments.

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